Porting optimized HPCG 3.1 for IBM BG/Q to IBM POWER9

Optimization overview and main results

Panagiotis Chatzidoukas, Cristiano Malossi, Christoph Hagleitner, Costas Bekas IBM Research – Zurich

November 13th, 2018



Analysis of HPCG 2.4 on BG/Q: no optimization

- System: Juqueen IBM BG/Q, 28 racks, 458752 cores
- HPCG local domain dimension: 80x80x80
- Result: 59.1 TFlop/s (2 GFlops/s per node)

Kernel	Time [%]	GFlop/s
DOT	1.5	62997.8
WAXPY	1.34	69182.5
SPMV	14.4	59170
MG	82.78	57495.1
Raw total	100	57971.4
Total	-	59154.4

Key observations:

- 82 % of time is MG (no threads in ref. implementation)
- MG main routine need to load 15 to 20 Bytes for each 2 Flops: Bytes/Flop = 7.5 to 10
- BG/Q has 30 GB/s bandwidth to memory, so we expect 3 to 4 GFlop/node

Optimizations (1) – Smart pivoting for parallel SYMGS

- BG/Q has 4 threads per core
- A parallel implementation of SYMGS requires coloring
- Coloring has two undesired side effects: 1) slow down convergence, 2) limit cache reuse

Our solution:

• Stencil discretization lead to a uniform (diagonal) matrix structure: we can rely on that!



Optimizations (2) – Fine tuning

- Strong compiler optimization (-O5, -qipa=level=2, -qhot=level=2, etc., little effect overall)
- AXPY and DOT manually SIMD vectorized (slightly better performances than auto-compiled versions)
- Contiguous storage for matrix (to help hardware prefetching very important on BG/Q)
- Improve backward prefetching of Gauss-Seidel smoother with __dcbt instructions
- Manual unrolling factor of 2 for SpMV, and slightly improved code
- Use Isend/Irecv with a single MPI_Waitall call at the end (better overlap of communications)

and, not less important, optimized local problem size

- A smaller problem size gives better MG and SPMV performance (better cache use)
 ... however it also increases DOT MPI_Allreduce (due to wait time) need to find the best compromise!
- We use 48x16x16 or 56x16x16 (on few racks) and 96x32x32 or 112x32x32 (on many racks)

HPCG 2.4 on BG/Q benchmark (reference vs optimized)

Non-optimized:

- 32 MPI task/node, no threads
- Local domain dimension: 80x80x80
- 59.1 Tflops/s (2 GFlop/s per node)

Kernel	Time [%]	GFlop/s
DOT	1.5	62997.8
WAXPY	1.34	69182.5
SPMV	14.4	59170
MG	82.78	57495.1
Raw total	100	57971.4
Total	-	59154.4

Optimized:

- 16 MPI task/node, 4 threads each
- Local domain dimension: 112x32x32
- 95.5 Tflop/s (3.33 GFlop/s per node)

Kernel	Time [%]	GFlop/s
DOT	3.68	42098.9
WAXPY	1.5	103352
SPMV	14.11	100137
MG	80.7	98033.7
Raw total	100	96333.2
Total	-	95476.4

5 | © 2018 - IBM Corporation Copyright | IBM Confidential

HPCG 3.1 on Vulcan & Sequoia

Vulcan (24K nodes):

6 | ©

- 32 MPI tasks/node, 2 threads each
- Local domain dimension: 112x32x32
- 80.9 TFlop/s (3.29 GFlop/s per node)

Sequoia (96K nodes):

- 32 MPI task/node, 2 threads each
- Local domain dimension: 112x32x32
- 330.4 Tflop/s (3.36 GFlop/s per node)
 #10 in HPGC results list (June 2018)

Kernel	Time [%]	GFlop/s	Kernel	Time [%]	GFlop/s
DOT	6.5	21420.3	DOT	3.8	149613
WAXPY	2.0	69445.9	WAXPY	1.8	324545
SPMV	14.2	88960.5	SPMV	15.3	343314
MG	77.3	91263.6	MG	79.1	370952
Raw total	100	85980	Raw total	100	357356
Total	-	80893.7 peak	= 1.6% Total	-	330373 peak =
18 - IBM Corporatio	on Copyright IBM Co	nfidential	perfect scaling fro	m 24K to 96K no	des

Changes in optimizations for POWER9

- The XLC (16.1.0) compiler optimizations for BGQ remain the same still little effect overall
- MPI configuration
 - One task per core
 - Binding policy: mpirun --bind-to core --map-by core
- OpenMP configuration:
 - two threads per task (core)
 - OMP_PROC_BIND=FALSE (no explicit binding to the hardware threads of the core)
 - OMP_WAIT_POLICY=ACTIVE (no need for yield)
- Problem size
 - Local domain is set to 160x160x96 (or 160x96x160)
 - Larger than BGQ due to higher memory bandwidth

MareNostrum P9 CTE @ BSC

- #255 in TOP500 (June 2018)
 - https://www.top500.org/system/179442
- 54 nodes organized in 3 racks
 - 52 compute and 2 login nodes
- Each node
 - 2x IBM POWER9 20C 3.1GHz
 - 40 cores with four-way multithreading variant SMT4
 - 4x NVIDIA Tesla V100
- Interconnection network: Dual-rail Mellanox EDR Infiniband

HPCG 3.1 on MareNostrum POWER9 CTE

Performance (GFlop/s)

200

100

MareNostrum IBM POWER9 CTE (16 nodes):

- 40 MPI tasks/node, 2 threads each
- Local domain dimension: 160x160x96
- 597 GFlop/s (37.3 GFlop/s per node)

Kernel	Time [%]	GFlop/s
DOT	5.8	173.3
WAXPY	2.7	374.4
SPMV	14.4	636.3
MG	77.1	666.3
Raw total	100	625.4
Total	-	597.4



8

Number of Nodes

4

12

16

HPGC on P9 CTE – CPU only

9 | © 2018 - IBM Corporation Copyright | IBM Confidential

HPCG 3.1 on MareNostrum POWER9 CTE

MareNostrum IBM POWER9 CTE (20 nodes):

- 40 MPI tasks/node, 2 threads each
- Local domain dimension: 160x160x96
- 744.6 GFlop/s (37.2 GFlop/s per node)

Kernel	Time [%]	GFlop/s
DOT	3.5	359.9
WAXPY	2.7	475.4
SPMV	14.9	771.0
MG	78.9	815.7
Raw total	100	783.9
Total	-	744.6



10 | © 2018 - IBM Corporation Copyright | IBM Confidential

HPCG 3.1 on MareNostrum POWER9 CTE

MareNostrum IBM POWER9 CTE (32 nodes):

- 40 MPI tasks/node, 2 threads each
- Local domain dimension: 160x160x96
- 1193.4 GFlop/s (37.3 GFlop/s per node)

Kernel	Time [%]	GFlop/s
DOT	6.5	311.8
WAXPY	2.7	740.4
SPMV	14.5	1268.7
MG	76.4	1342.8
Raw total	100	1249.0
Total	-	1193.4





Conclusions

- An optimized CPU-only version of HPCG on IBM processors
 - explicit SIMD vectorization, data prefetching, asynchronous MPI communication
 - smart pivoting: new OpenMP parallelization approach for SYMGS, the most time consuming kernel of HPCG
- Fine tuning of several parameters, such as
 - Local problem size
 - Number of MPI tasks and OpenMP threads
- The code achieves 1.6% of the peak performance on IBM BGQ and 3.9% on IBM POWER9
- The code will be released soon as open source project

THANK YOU

IBM, the IBM logo, and ibm.com are trademarks or registered trademarks of International Business Machines Corporation in the United States, other countries, or both. Other product and service names might be trademarks of IBM or other companies.



